Speech and Language Processing

Parsing
Chapter 13

Today

- Parsing with CFGs
 - Bottom-up, top-down
 - Ambiguity
 - CKY parsing
 - Earley parsing

Parsing

- Parsing with CFGs refers to the task of assigning proper trees to input strings
- Proper here means a tree that covers all and only the elements of the input and has an S at the top
- It doesn't actually mean that the system can select the correct tree from among all the possible trees

Parsing

- As with everything of interest, parsing involves a search which involves the making of choices
- We'll start with some basic (meaning bad) methods before moving on to the one or two that you need to know

Sample L₁ Grammar

Grammar	Lexicon
$S \rightarrow NP VP$	$Det \rightarrow that \mid this \mid a$
$S \rightarrow Aux NP VP$	$Noun \rightarrow book \mid flight \mid meal \mid money$
$S \rightarrow VP$	$Verb \rightarrow book \mid include \mid prefer$
$NP \rightarrow Pronoun$	$Pronoun \rightarrow I \mid she \mid me$
$NP \rightarrow Proper-Noun$	Proper-Noun → Houston NWA
$NP \rightarrow Det\ Nominal$	$Aux \rightarrow does$
$Nominal \rightarrow Noun$	$Preposition \rightarrow from \mid to \mid on \mid near \mid through$
$Nominal \rightarrow Nominal Noun$	
$Nominal \rightarrow Nominal PP$	
$VP \rightarrow Verb$	
$VP \rightarrow Verb NP$	
$VP \rightarrow Verb NP PP$	
$VP \rightarrow Verb PP$	
$VP \rightarrow VP PP$	
$PP \rightarrow Preposition NP$	

For Now

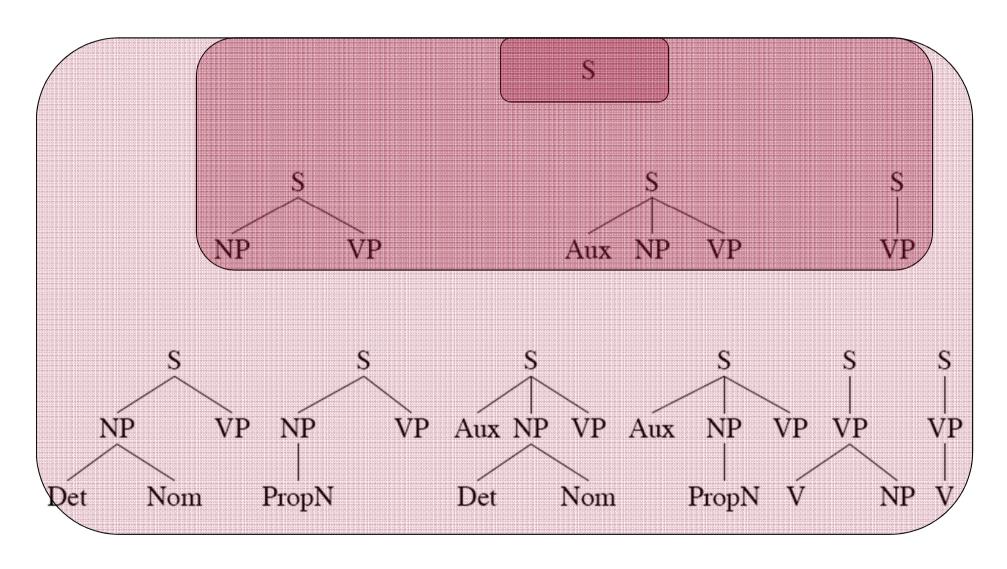
Assume...

- You have all the words already in some buffer
- The input isn't POS tagged
- We won't worry about morphological analysis
- All the words are known
- These are all problematic in various ways, and would have to be addressed in real applications.

Top-Down Search

- Since we're trying to find trees rooted with an S (Sentences), why not start with the rules that give us an S.
- Then we can work our way down from there to the words.
- Example sentence: "Book that flight"

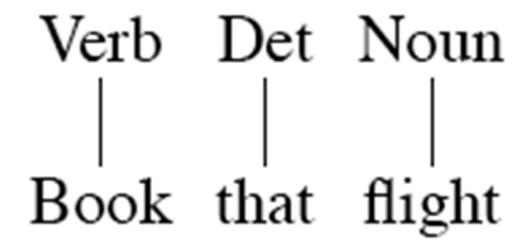
Top Down Space

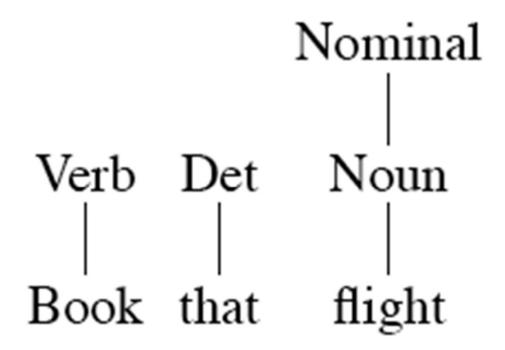


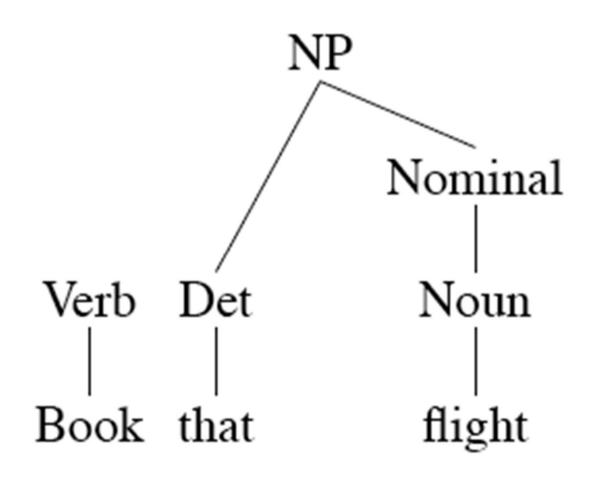
Bottom-Up Parsing

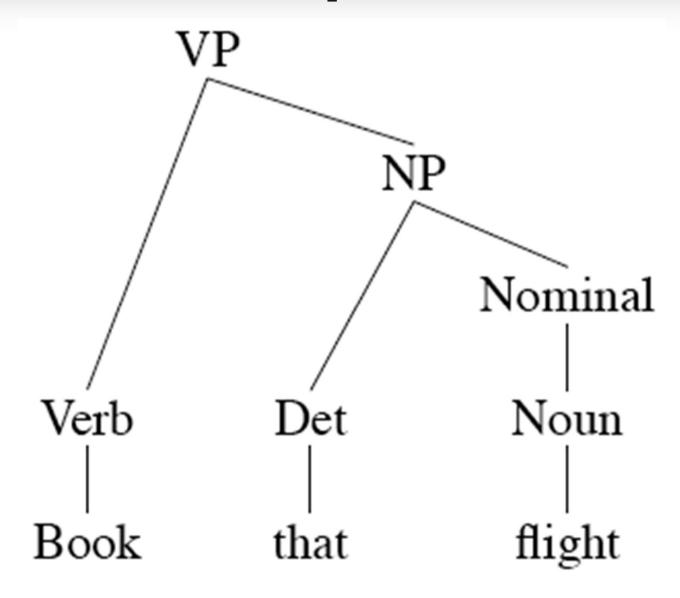
- Of course, we also want trees that cover the input words. So we might also start with trees that link up with the words in the right way.
- Then work your way up from there to larger and larger trees.

Book that flight









Top-Down and Bottom-Up

Top-down

- Only searches for trees that can be answers (i.e. S's)
- But also suggests trees that are not consistent with any of the words

Bottom-up

- Only forms trees consistent with the words
- But suggests trees that make no sense globally

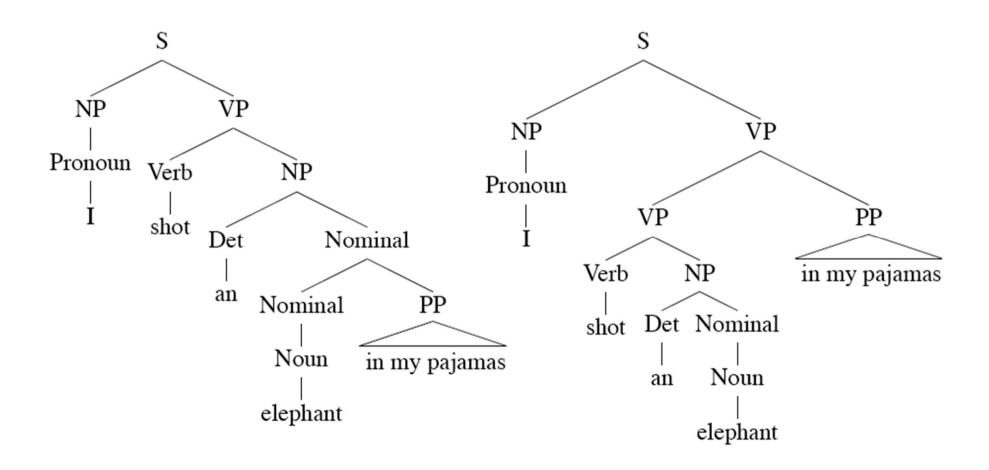
Control

- Of course, in both cases we left out how to keep track of the search space and how to make choices
 - Which node to try to expand next
 - Which grammar rule to use to expand a node
- One approach is called backtracking.
 - Make a choice, if it works out then fine
 - If not then back up and make a different choice

Problems

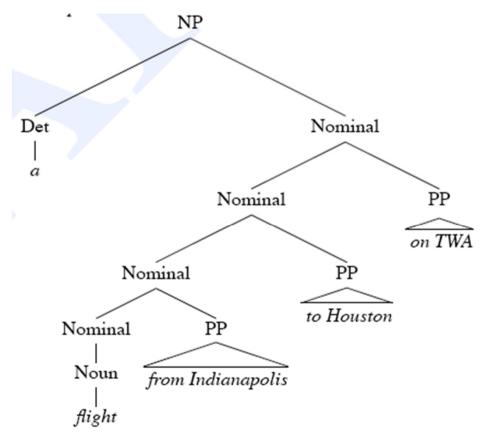
- Even with the best filtering, backtracking methods are doomed because of two inter-related problems
 - Ambiguity
 - Shared subproblems

Ambiguity

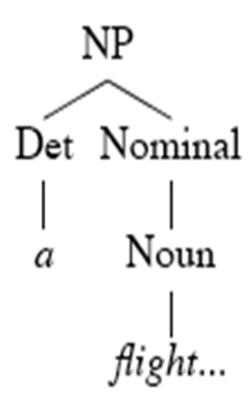


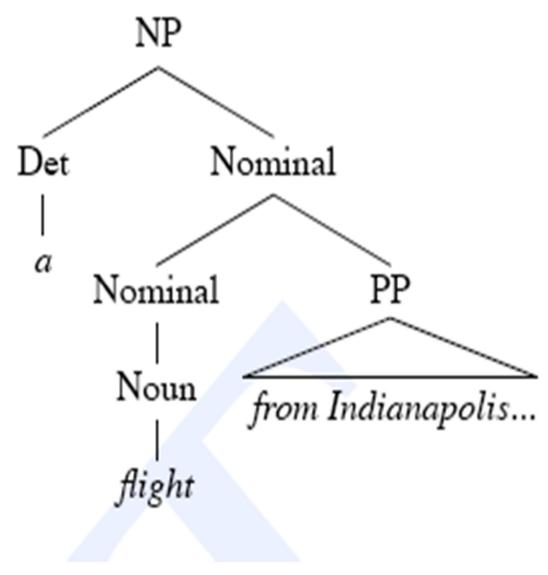
- No matter what kind of search (top-down or bottom-up or mixed) that we choose.
 - We don't want to redo work we've already done.
 - Unfortunately, naïve backtracking will lead to duplicated work.

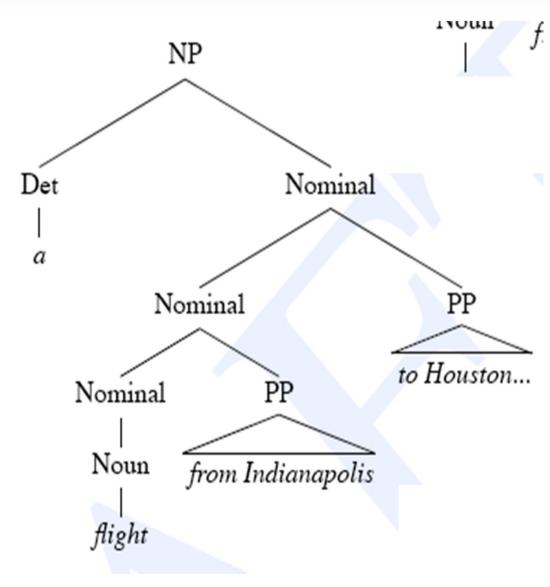
- Consider
 - A flight from Indianapolis to Houston on TWA

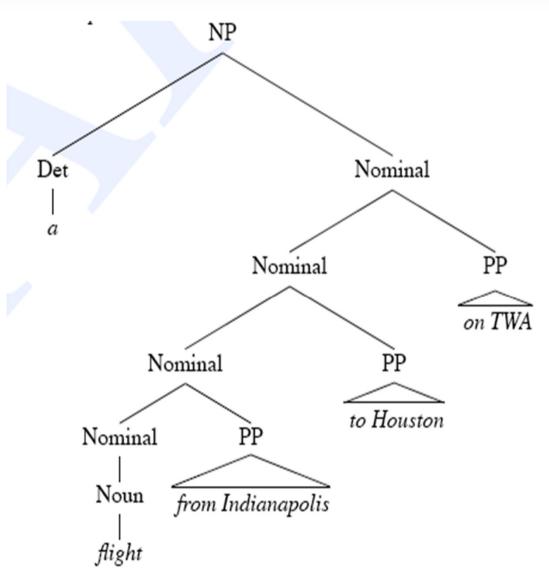


- Assume a top-down parse making choices among the various Nominal rules.
- In particular, between these two
 - Nominal -> Noun
 - Nominal -> Nominal PP
- Statically choosing the rules in this order leads to the following bad results...









Dynamic Programming

- DP search methods fill tables with partial results and thereby
 - Avoid doing avoidable repeated work
 - Solve exponential problems in polynomial time Efficiently store ambiguous structures with shared sub-parts.
- We'll cover two approaches that roughly correspond to top-down and bottom-up approaches.
 - CKY (bottom-up)
 - Earley (top-down)

CKY Parsing

- First we'll limit our grammar to epsilonfree, binary rules (more later)
- Consider the rule $A \rightarrow BC$
 - If there is an A somewhere in the input then there must be a B followed by a C in the input.
 - If the A spans from i to j in the input then there must be some k such that i<k<j</p>
 - Ie. The B splits from the C someplace.

Problem

- What if your grammar isn't binary?
 - As in the case of the TreeBank grammar?
- Convert it to binary... any arbitrary CFG can be rewritten into Chomsky-Normal Form automatically.
- What does this mean?
 - The resulting grammar accepts (and rejects) the same set of strings as the original grammar.
 - But the resulting derivations (trees) are different.

Problem

 More specifically, we want our rules to be of the form

$$A \rightarrow B C$$

Or

$$A \rightarrow W$$

That is, rules can expand to either 2 nonterminals or to a single terminal.

Binarization Intuition

- Eliminate chains of unit productions.
- Introduce new intermediate non-terminals into the grammar that distribute rules with length > 2 over several rules.
 - So... $S \rightarrow A B C turns into$

$$S \rightarrow X C$$
 and

$$X \rightarrow A B$$

Where X is a symbol that doesn't occur anywhere else in the the grammar.

Sample L1 Grammar

Grammar	Lexicon
$S \rightarrow NP VP$	$Det \rightarrow that \mid this \mid a$
$S \rightarrow Aux NP VP$	$Noun \rightarrow book \mid flight \mid meal \mid money$
$S \rightarrow VP$	$Verb \rightarrow book \mid include \mid prefer$
$NP \rightarrow Pronoun$	$Pronoun \rightarrow I \mid she \mid me$
$NP \rightarrow Proper-Noun$	Proper-Noun $ ightarrow$ Houston NWA
$NP \rightarrow Det Nominal$	$Aux \rightarrow does$
$Nominal \rightarrow Noun$	$Preposition \rightarrow from \mid to \mid on \mid near \mid through$
$Nominal \rightarrow Nominal Noun$	
$Nominal \rightarrow Nominal PP$	
$VP \rightarrow Verb$	
$VP \rightarrow Verb NP$	
$VP \rightarrow Verb NP PP$	
$VP \rightarrow Verb PP$	
$VP \rightarrow VP PP$	
$PP \rightarrow Preposition NP$	

CNF Conversion

\mathscr{L}_1 Grammar	\mathscr{L}_1 in CNF
$S \rightarrow NP VP$	$S \rightarrow NP VP$
$S \rightarrow Aux NP VP$	$S \rightarrow X1 VP$
	$X1 \rightarrow Aux NP$
$S \rightarrow VP$	$S \rightarrow book \mid include \mid prefer$
	$S \rightarrow Verb NP$
	$S \rightarrow X2 PP$
	$S \rightarrow Verb PP$
	$S \rightarrow VPPP$
$NP \rightarrow Pronoun$	$NP \rightarrow I \mid she \mid me$
$NP \rightarrow Proper-Noun$	$NP \rightarrow TWA \mid Houston$
$NP \rightarrow Det\ Nominal$	$NP \rightarrow Det Nominal$
Nominal → Noun	$Nominal \rightarrow book \mid flight \mid meal \mid money$
Nominal → Nominal Noun	$Nominal \rightarrow Nominal Noun$
$Nominal \rightarrow Nominal PP$	$Nominal \rightarrow Nominal PP$
$VP \rightarrow Verb$	$VP \rightarrow book \mid include \mid prefer$
$VP \rightarrow Verb NP$	$VP \rightarrow Verb NP$
$VP \rightarrow Verb NP PP$	$VP \rightarrow X2 PP$
	$X2 \rightarrow Verb NP$
$VP \rightarrow Verb PP$	$VP \rightarrow Verb PP$
$VP \rightarrow VP PP$	$VP \rightarrow VP PP$
$PP \rightarrow Preposition NP$	$PP \rightarrow Preposition NP$

CKY

- So let's build a table so that an A spanning from i to j in the input is placed in cell [i,j] in the table.
- So a non-terminal spanning an entire string will sit in cell [0, n]
 - Hopefully an S
- If we build the table bottom-up, we'll know that the parts of the A must go from i to k and from k to j, for some k.

CKY

- Meaning that for a rule like A → B C we should look for a B in [i,k] and a C in [k,j].
- In other words, if we think there might be an A spanning i,j in the input... AND
 - $A \rightarrow B C$ is a rule in the grammar THEN
- There must be a B in [i,k] and a C in [k,j] for some i<k<j</p>

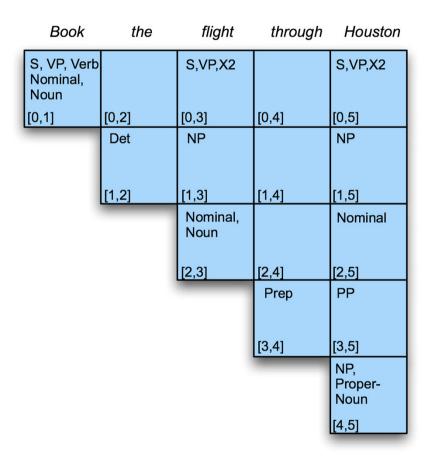
CKY

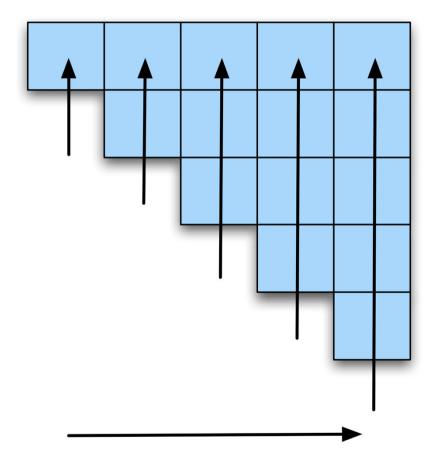
- So to fill the table, loop over the cell[i,j] values in some systematic way
 - What constraint should we put on that systematic search?
 - For each cell, loop over the appropriate k values to search for things to add.

CKY Algorithm

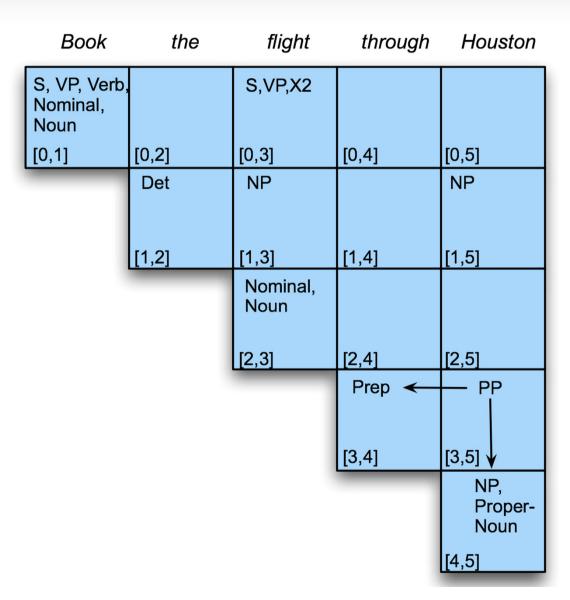
function CKY-PARSE(words, grammar) returns table

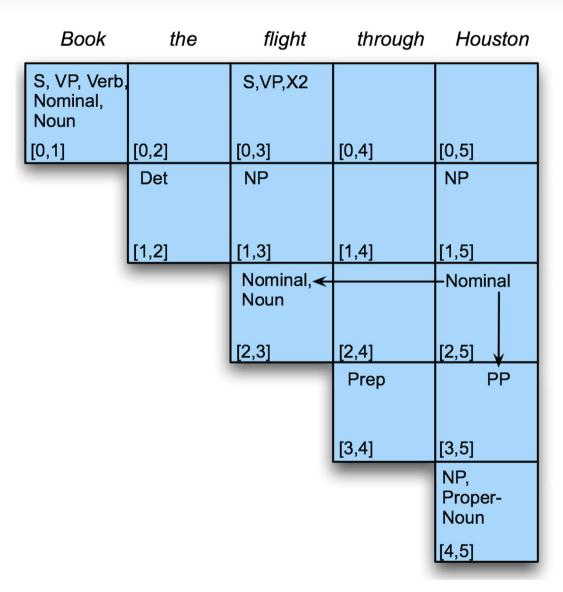
```
for j \leftarrow from 1 to LENGTH(words) do table[j-1,j] \leftarrow \{A \mid A \rightarrow words[j] \in grammar\} for i \leftarrow from j-2 downto 0 do table[i,j] \leftarrow table[i,j] \cup \{A \mid A \rightarrow BC \in grammar, B \in table[i,k], C \in table[k,j]\}
```



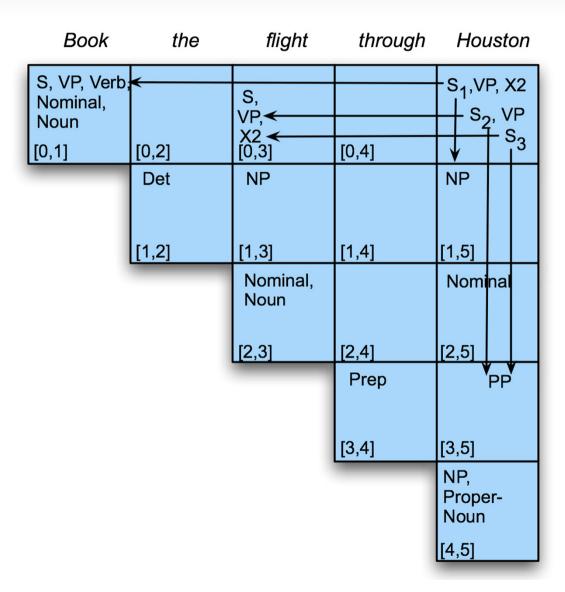


	Book	the	flight	through	Houston
	S, VP, Verb, Nominal, Noun		S,VP,X2		
	[0,1]	[0,2]	[0,3]	[0,4]	[0,5]
		Det	NP		
		[1,2]	[1,3]	[1,4]	[1,5]
			Nominal, Noun		Nominal
			[2,3]	[2,4]	[2,5]
				Prep	
Filling colum	n 5			[3,4]	[3,5]
<u> </u>					NP, Proper- Noun
					[4,5]





Book	the	flight	through	Houston
S, VP, Verb, Nominal, Noun		S,VP,X2		
[0,1]	[0,2]	[0,3]	[0,4]	[0,5]
	Det ←	NP		NP
	[1,2]	[1,3]	[1,4]	[15]
		Nominal, Noun		Nominal
		[2,3]	[2,4]	[2,5]
			Prep	PP
			[3,4]	[3,5]
				NP, Proper- Noun
				[4,5]



CKY Parsing

- Is that really a parser?
- We need to augment the entries in the table so that each non-terminal is paired with pointers to the table entries from which it was derived.

Note

- We arranged the loops to fill the table a column at a time, from left to right, bottom to top.
 - This assures us that whenever we're filling a cell, the parts needed to fill it are already in the table (to the left and below)
 - It's somewhat natural in that it processes the input a left to right a word at a time
 - Known as online

CKY Notes

- Since it's bottom up, CKY populates the table with a lot of phantom constituents.
 - Segments that by themselves are constituents but cannot really occur in the context in which they are being suggested.
 - To avoid this we can switch to a top-down control strategy
 - Or we can add some kind of filtering that blocks constituents where they can not happen in a final analysis.

Earley Parsing

- Allows arbitrary CFGs
- Top-down control
- Fills a table in a single sweep over the input
 - Table is length N+1; N is number of words
 - Table entries represent
 - Completed constituents and their locations
 - In-progress constituents
 - Predicted constituents

States

 The table-entries are called states and are represented with dotted-rules.

 $S \rightarrow VP$

A VP is predicted

NP → Det * Nominal

An NP is in progress

 $VP \rightarrow V NP$

A VP has been found (completed)

States/Locations

$$\blacksquare$$
 S $\rightarrow \bullet$ VP [0,0]

 A VP is predicted at the start of the sentence

 An NP is in progress; the Det goes from 1 to 2

• $VP \rightarrow V NP \bullet [0,3]$

 A VP has been found starting at 0 and ending at 3

Earley

- As with most dynamic programming approaches, the answer is found by looking in the table in the right place.
- In this case, there should be an S state in the final column that spans from 0 to N and is complete. That is,
 - \blacksquare S $\rightarrow \alpha \bullet [0,N]$
- If that's the case, you're done.

Earley

- So sweep through the table from 0 to N...
 - New **predicted** states are created by starting top-down from S
 - New incomplete states are created by advancing existing states as new constituents are discovered
 - New complete states are created in the same way.

Earley

- More specifically...
 - 1. Predict all the states you can upfront
 - 2. Read a word
 - 1. Extend states based on matches
 - 2. Generate new predictions
 - 3. Go to step 2
 - 3. When you're out of words, look at the chart to see if you have a winner

Core Earley Code

function EARLEY-PARSE(words, grammar) returns chart

```
ENQUEUE((\gamma \rightarrow \bullet S, [0,0]), chart[0])
for i \leftarrow from 0 to LENGTH(words) do
 for each state in chart[i] do
   if INCOMPLETE?(state) and
            NEXT-CAT(state) is not a part of speech then
      Predictor(state)
   elseif INCOMPLETE?(state) and
            NEXT-CAT(state) is a part of speech then
       SCANNER(state)
   else
      COMPLETER(state)
 end
end
return(chart)
```

Earley Code

```
procedure PREDICTOR((A \rightarrow \alpha \bullet B \beta, [i, j]))
   for each (B \rightarrow \gamma) in GRAMMAR-RULES-FOR(B, grammar) do
         ENQUEUE((B \rightarrow \bullet \gamma, [j, j]), chart[j])
   end
procedure SCANNER((A \rightarrow \alpha \bullet B \beta, [i, j]))
   if B \subset PARTS-OF-SPEECH(word[j]) then
       ENQUEUE((B \rightarrow word[j], [j, j+1]), chart[j+1])
procedure COMPLETER((B \rightarrow \gamma \bullet, [j,k]))
   for each (A \rightarrow \alpha \bullet B \beta, [i, j]) in chart[j] do
         ENQUEUE((A \rightarrow \alpha B \bullet \beta, [i,k]), chart[k])
    end
```

- Book that flight
- We should find... an S from 0 to 3 that is a completed state...

Chart[0]

S0	$\gamma \rightarrow \bullet S$	[0,0]	Dummy start state
S1	$S \rightarrow \bullet NP VP$	[0,0]	Predictor
S2	$S \rightarrow \bullet Aux NP VP$	[0,0]	Predictor
S3	$S \rightarrow \bullet VP$	[0,0]	Predictor
S4	$NP \rightarrow \bullet Pronoun$	[0,0]	Predictor
S5	$NP \rightarrow \bullet Proper-Noun$	[0,0]	Predictor
S6	NP ightarrow ullet Det Nominal	[0,0]	Predictor
S7	$VP \rightarrow \bullet Verb$	[0,0]	Predictor
S8	$VP \rightarrow \bullet Verb NP$	[0,0]	Predictor
S9	$VP \rightarrow \bullet Verb NP PP$	[0,0]	Predictor
S10	$VP \rightarrow \bullet Verb PP$	[0,0]	Predictor
S11	$VP \rightarrow \bullet VP PP$	[0,0]	Predictor

Note that given a grammar, these entries are the same for all inputs; they can be pre-loaded.

Chart[1]

S12	$Verb \rightarrow book \bullet$	[0,1]	Scanner
S13	$VP \rightarrow Verb \bullet$	[0,1]	Completer
S14	$VP \rightarrow Verb \bullet NP$	[0,1]	Completer
S15	$VP \rightarrow Verb \bullet NP PP$	[0,1]	Completer
S16	$VP \rightarrow Verb \bullet PP$	[0,1]	Completer
S17	$S \rightarrow VP \bullet$	[0,1]	Completer
S18	$VP \rightarrow VP \bullet PP$	[0,1]	Completer
S19	$NP \rightarrow \bullet Pronoun$	[1,1]	Predictor
S20	$NP \rightarrow \bullet Proper-Noun$	[1,1]	Predictor
S21	$NP \rightarrow \bullet Det Nominal$	[1,1]	Predictor
S22	$PP \rightarrow \bullet Prep NP$	[1,1]	Predictor

Charts[2] and [3]

$Det \rightarrow that \bullet$	[1,2]	Scanner
$NP \rightarrow Det \bullet Nominal$	[1,2]	Completer
$Nominal \rightarrow \bullet Noun$	[2,2]	Predictor
$Nominal \rightarrow \bullet Nominal Noun$	[2,2]	Predictor
$Nominal \rightarrow \bullet Nominal PP$	[2,2]	Predictor
$Noun \rightarrow flight \bullet$	[2,3]	Scanner
$Nominal \rightarrow Noun \bullet$	[2,3]	Completer
NP o Det Nominal ullet	[1,3]	Completer
$Nominal \rightarrow Nominal \bullet Noun$	[2,3]	Completer
$Nominal \rightarrow Nominal \bullet PP$	[2,3]	Completer
$VP \rightarrow Verb NP \bullet$	[0,3]	Completer
$VP \rightarrow Verb NP \bullet PP$	[0,3]	Completer
$PP \rightarrow \bullet Prep NP$	[3,3]	Predictor
$S \rightarrow VP \bullet$	[0,3]	Completer
$VP \rightarrow VP \bullet PP$	[0,3]	Completer
	$NP o Det ullet Nominal \ Nominal o ullet Noun \ Nominal o ullet Nominal Noun \ Nominal o ullet Nominal PP \ Noun o flight ullet \ Nominal o Noun ullet \ Nominal o Noun ullet \ Nominal o Nominal ullet Nominal ullet Noun \ Nominal o Nominal ullet Nominal ullet PP \ VP o Verb NP ullet \ VP o Verb NP ullet PP \ PP o ullet Prep NP \ S o VP ullet \ VP ullet \ VP ullet \ VP ullet \ Prep NP \ S o VP ullet \ PP \ Prep NP \ S o VP ullet \ PP \ PP \ PP \ PP \ Prep NP \ Prep NP \ PP $	$NP o Det ullet Nominal \qquad [1,2]$ $Nominal o ullet Noun \qquad [2,2]$ $Nominal o ullet Nominal Noun \qquad [2,2]$ $Nominal o ullet Nominal PP \qquad [2,2]$ $Noun o flight ullet \qquad [2,3]$ $Nominal o Noun ullet \qquad [2,3]$ $NP o Det Nominal ullet \qquad [1,3]$ $Nominal o Nominal ullet Noun \qquad [2,3]$ $Nominal o Nominal ullet Nominal ullet PP \qquad [2,3]$ $VP o Verb \ NP ullet \qquad [0,3]$ $VP o Verb \ NP ullet PP \qquad [0,3]$ $PP o ullet Prep \ NP \qquad [3,3]$ $S o VP ullet \qquad [0,3]$

Efficiency

- For such a simple example, there seems to be a lot of useless stuff in there.
- Why?

- It's predicting things that aren't consistent with the input
- ·That's the flipside to the CKY problem.

Details

- As with CKY that isn't a parser until we add the back-pointers so that each state knows where it came from
 - Add an additional field to each state with information about the completed states that generated its constituents.

Chart[1]	S12	Verb->book •	[0,1]	Scanner
Chart[2]	S23	Det->that •	[1,2]	Scanner
Chart[3]	S28	Noun->flight •	[2,3]	Scanner
	S29	Nominal->Noun •	[2,3]	(S28)
	S30	NP -> Det Nominal •	[1,3]	(S23, S29)
	S33	VP -> Verb NP •	[0,3]	(S12, S30)
	S36	S -> VP •	[0,3]	(S33)

Back to Ambiguity

Did we solve it?

Ambiguity

No...

- Both CKY and Earley will result in multiple S structures for the [0,N] table entry.
- They both efficiently store the sub-parts that are shared between multiple parses.
- And they obviously avoid re-deriving those sub-parts.
- But neither can tell us which one is right.

Ambiguity

- In most cases, humans don't notice incidental ambiguity (lexical or syntactic). It is resolved on the fly and never noticed.
- We'll try to model that with probabilities.