

First-Order Logic

Chapter 8

Outline

- Why FOL?
- Syntax and semantics of FOL
- Using FOL
- Wumpus world in FOL
- Knowledge engineering in FOL

Propositional logic, pros and cons

- ☺ Propositional logic is **declarative**
- ☺ Propositional logic allows partial (disjunctive/negated) information
 - (unlike most data structures and databases)
- ☺ Propositional logic is **compositional**:
 - meaning of $B_{1,1} \wedge P_{1,2}$ is derived from meaning of $B_{1,1}$ and of $P_{1,2}$

Propositional logic, pros and cons

- ⊙ Meaning in propositional logic is **context-independent**
 - (unlike natural language, where meaning depends on context)
- ⊙ Propositional logic has very limited expressive power
 - (unlike natural language)
 - E.g., cannot say "pits cause breezes in adjacent squares"
 - except by writing one sentence for each square

Why not use Natural Language?

- It serves a different purpose:
 - Communication rather than representation
- It is not compositional
 - Context matters
- It can be ambiguous
 - Again, context matters

Create a new language

- Builds on propositional logic
- But is inspired by natural language!

First-order logic

- Whereas propositional logic assumes the world contains **facts**,
- first-order logic (like natural language) assumes the world contains
 - **Objects**: people, houses, numbers, colors, baseball games, wars, ...
 - **Relations**: red, round, prime, brother of, bigger than, part of, comes between, ...
 - **Functions**: father of, best friend, one more than, plus, ...

Syntax of FOL: Basic elements

- **Constants** KingJohn, 2, NUS,...
- **Predicates** Brother, >,...
- **Functions** Sqrt, LeftLegOf,...
- **Variables** x, y, a, b,...
- **Connectives** $\neg, \Rightarrow, \wedge, \vee, \Leftrightarrow$
- **Equality** =
- **Quantifiers** \forall, \exists

Atomic sentences

Atomic sentence = $predicate(term_1, \dots, term_n)$
or $term_1 = term_2$

Term = $function(term_1, \dots, term_n)$
or *constant* or *variable*

- E.g., $Brother(KingJohn, RichardTheLionheart)$
 $>(Length(LeftLegOf(Richard)), Length(LeftLegOf(KingJohn)))$

Complex sentences

- Complex sentences are made from atomic sentences using connectives

$$\neg S, S_1 \wedge S_2, S_1 \vee S_2, S_1 \Rightarrow S_2, S_1 \Leftrightarrow S_2$$

E.g. $Sibling(KingJohn, Richard) \Rightarrow Sibling(Richard, KingJohn)$

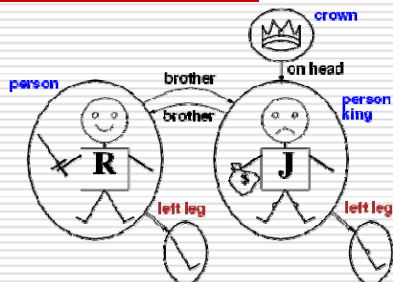
$$>(1,2) \vee \leq(1,2)$$

$$<(1,2) \wedge \neg >(1,2)$$

Truth in first-order logic

- Sentences are true with respect to a **model** and an **interpretation (DIAGRAM)**
- Model contains objects (**domain elements**) and relations among them
- Interpretation specifies referents for
 - constant symbols \rightarrow objects
 - predicate symbols \rightarrow relations
 - function symbols \rightarrow functional relations
- An atomic sentence $predicate(term_1, \dots, term_n)$ is true iff the **objects** referred to by $term_1, \dots, term_n$ are in the **relation** referred to by $predicate$

Models for FOL: Example



Universal quantification

- $\forall <variables> <sentence>$

Everyone in HR is smart:

$$\forall x \text{ At}(x, \text{HR}) \Rightarrow \text{Smart}(x)$$

- $\forall x P$ is true in a model m iff P is true with x being each possible object in the model
- Roughly speaking, equivalent to the conjunction of instantiations of P
 - At(KingJohn,HR) \Rightarrow Smart(KingJohn)
 - \wedge At(Richard,HR) \Rightarrow Smart(Richard)
 - \wedge At(HR,HR) \Rightarrow Smart(HR)
 - $\wedge \dots$

A common mistake to avoid

- Typically, \Rightarrow is the main connective with \forall
- Common mistake: using \wedge as the main connective with \forall :
 $\forall x \text{ At}(x, \text{HR}) \wedge \text{Smart}(x)$
means "Everyone is at HR and everyone is smart"

Existential quantification

- $\exists <variables> <sentence>$

Someone at HR is smart:

$$\exists x \text{ At}(x, \text{HR}) \wedge \text{Smart}(x)$$

- $\exists x P$ is true in a model m iff P is true with x being some possible object in the model
- Roughly speaking, equivalent to the disjunction of instantiations of P
 - At(KingJohn,HR) \wedge Smart(KingJohn)
 - \vee At(Richard,HR) \wedge Smart(Richard)
 - \vee At(HR,HR) \wedge Smart(HR)
 - $\vee \dots$

Another mistake to avoid

□ Typically, \wedge is the main connective with \exists

□ Common mistake: using \Rightarrow as the main connective with \exists :

$$\exists x \text{ At}(x, \text{HR}) \Rightarrow \text{Smart}(x)$$

is true if there is anyone who is not at HR!

Properties of quantifiers

□ $\forall x \forall y$ is the same as $\forall y \forall x$

□ $\exists x \exists y$ is the same as $\exists y \exists x$

□ $\exists x \forall y$ is **not** the same as $\forall y \exists x$

□ $\exists x \forall y \text{ Loves}(x, y)$

■ "There is a person who loves everyone in the world"

□ $\forall y \exists x \text{ Loves}(x, y)$

■ "Everyone in the world is loved by at least one person"

□ **Quantifier duality**: each can be expressed using the other

$$\forall x \text{ Likes}(x, \text{IceCream}) \quad \neg \exists x \neg \text{Likes}(x, \text{IceCream})$$

$$\exists x \text{ Likes}(x, \text{Broccoli}) \quad \neg \forall x \neg \text{Likes}(x, \text{Broccoli})$$

Equality

□ $term_1 = term_2$ is true under a given interpretation if and only if $term_1$ and $term_2$ refer to the same object

□ E.g., definition of *Sibling* in terms of *Parent*:

$$\forall x, y \text{ Sibling}(x, y) \Leftrightarrow$$

$$[\neg(x = y) \wedge \exists m, f \neg (m = f) \wedge \text{Parent}(m, x) \wedge \text{Parent}(f, x) \wedge \text{Parent}(m, y) \wedge \text{Parent}(f, y)]$$

Using FOL

The kinship domain:

- Brothers are siblings
 $\forall x,y \text{ Brother}(x,y) \Leftrightarrow \text{Sibling}(x,y)$
- One's mother is one's female parent
 $\forall m,c \text{ Mother}(c) = m \Leftrightarrow (\text{Female}(m) \wedge \text{Parent}(m,c))$
- "Sibling" is symmetric
 $\forall x,y \text{ Sibling}(x,y) \Leftrightarrow \text{Sibling}(y,x)$

Some sentences are **Axioms** (i.e. definitions, facts) while others are **Theorems** derived from those.

Interacting with FOL KBs

- Suppose a wumpus-world agent is using an FOL KB and perceives a smell and a breeze (but no glitter) at $t=5$:

```
Tell(KB,Percept([Smell,Breeze,None],5))
Ask(KB,∃a BestAction(a,5))
```

- I.e., does the KB entail some best action at $t=5$?

- Answer: *Yes, {a/Shoot}* ← substitution (binding list)

- Given a sentence S and a substitution q ,
- Sq denotes the result of plugging q into S ; e.g.,
 $S = \text{Smarter}(x,y)$
 $q = \{x/\text{Hillary}, y/\text{Bill}\}$
 $Sq = \text{Smarter}(\text{Hillary}, \text{Bill})$

- $\text{Ask}(\text{KB}, S)$ returns some/all q such that $\text{KB} \models Sq$

KB for the wumpus world

- Perception
 - $\forall t,s,b \text{ Percept}([s,b,\text{Glitter}],t) \Rightarrow \text{Glitter}(t)$

- Reflex
 - $\forall t \text{ Glitter}(t) \Rightarrow \text{BestAction}(\text{Grab},t)$

Deducing hidden properties

- $\forall x,y,a,b \text{ Adjacent}([x,y],[a,b]) \Leftrightarrow [a,b] \in \{[x+1,y], [x-1,y],[x,y+1],[x,y-1]\}$

Properties of squares:

- $\forall s,t \text{ At}(\text{Agent},s,t) \wedge \text{Breeze}(t) \Rightarrow \text{Breezy}(s)$

Squares are breezy near a pit:

- **Diagnostic** rule---infer cause from effect
 $\forall s \text{ Breezy}(s) \Rightarrow \exists r \text{ Adjacent}(r,s) \wedge \text{Pit}(r)$
- **Causal** rule---infer effect from cause
 $\forall r \text{ Pit}(r) \Rightarrow [\forall s \text{ Adjacent}(r,s) \Rightarrow \text{Breezy}(s)]$

Summary

- First-order logic:
 - **objects** and **relations** are semantic primitives
 - syntax: constants, functions, predicates, equality, quantifiers
- Increased expressive power:
sufficient to define wumpus world
